## Learning DFA Decompositions from Examples and Demonstrations



Niklas Lauffer*, Beyazit Yalcinkaya*, Marcell Vazquez-Chanlatte, Ameesh Shah, and Sanjit A. Seshia
*equal contribution
University of California, Berkeley

## A simple gridworld



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$$
\text { Actions }=\{\wedge \vee \forall H\}
$$

## A simple gridworld



Actions $=\{\wedge \forall<\rightarrow\}$
$\operatorname{Pr}(\operatorname{slip} \downarrow)=1 / 32$

## A simple gridworld



$$
\operatorname{Pr}(\text { slip } \downarrow)=1 / 32
$$

## Specification mining



## Learning from demonstrations



## Learning decompositions



## Monolithic specifications can often be difficult to understand



## System-level specifications are often conjunctions of sub-specifications



## Inductive bias matters when learning from few demonstrations



## Contributions

1. SAT-based encoding for identifying a DFA decomposition of a specific size from labeled examples

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2. An algorithm for enumerating the full Pareto-frontier of decompositions

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2. An algorithm for enumerating the full Pareto-frontier of decompositions
3. Experimental analysis and extension to learning from demonstrations

# Structure of the talk 

\author{

1. Technical details <br> 2. Scalability analysis <br> 3. Learning from demonstrations
}

## State merging via coloring



Positive: \{a\} Negative: $\{\varnothing, \mathrm{b}, \mathrm{ba}, \mathrm{ab}, \mathrm{aab}\}$

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Positive: \{a\} Negative: $\{\varnothing, \mathrm{b}, \mathrm{ba}, \mathrm{ab}, \mathrm{aab}\}$

## State merging via coloring



Positive: \{a\} Negative: \{Ø, b, ba, ab, aab\}

## State merging for decompositions



Positive: $\{a\} \quad$ Negative: $\{\varnothing, b, b a, a b, a a b\}$

## State merging for decompositions



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## State merging for decompositions



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## A SAT encoding

Implemented as an extension of existing work*
*Ulyantsev, Vladimir \& Zakirzyanov, Ilya \& Shalyto, Anatoly. (2015). BFS-based Symmetry Breaking Predicates for DFA Identification

## A SAT encoding

Implemented as an extension of existing work*

- Each negative example must be rejected by at least one DFA:

$$
\bigwedge_{v \in V_{-}} \bigvee_{k \in[n]} \bigwedge_{i \in\left[m_{k}\right]} x_{v, i}^{k} \Longrightarrow \neg z_{i}^{k} .
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- Accepting and rejecting states of individual prefix trees cannot be merged:



## Finding the Pareto front of minimality



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## Structure of the talk

| Examples |  |
| :---: | :---: |
| Positive | Negative |
| a | $\emptyset$ |
|  | b |
|  | ba |
|  | ab |
|  | aab |

1. Technical details


## 2. Scalability analysis

3. Learning from demonstrations

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## Overhead comparable to the monolithic baseline

- Baseline, 2 Symbols, 4 DFAs, Time
- This Work, 2 Symbols, 4 DFAs, Time
$\square$ Baseline, 2 Symbols, 4 DFAs, Time CountThis Work, 2 Symbols, 4 DFAs, Time Count



## Overhead comparable to the monolithic baseline

- Baseline, 2 Symbols, 4 DFAs, Time
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= = This Work, 4 Symbols, 2 DFAs, TimeBaseline, 2 Symbols, 4 DFAs, Time Count Baseline, 4 Symbols, 2 DFAs, Time Count This Work, 2 Symbols, 4 DFAs, Time Count This Work, 4 Symbols, 2 DFAs, Time Count



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## Learning from demonstrations


$\operatorname{Pr}(\operatorname{slip} \downarrow)=1 / 32$

# Demonstration Informed Specification Search (DISS) 

## Learning from Demonstrations

## Demonstration Informed Specification Search (DISS)



## A helpful inductive bias from decompositions


 ever touch ㅆㅆ, you must then touch $\square$ before reaching ${ }^{[3}$.

## A helpful inductive bias from decompositions



Identified monolithic DFA (incorrect)
Reach $\overline{3}$ while avoiding $\mathbb{1}$. If you ever touch $\times$, you must then touch $\square$ before reaching ${ }^{[3}$.

## A helpful inductive bias from decompositions



Reach 3 while avoiding $\mathbb{C}$. If you ever touch $\times$, you must then touch $\square$ before reaching $\boldsymbol{\pi}$.


Identified monolithic DFA (incorrect)


Identified DFA decomposition (correct)

## Conclusion



- Known symmetry-breaking optimization still missing from the encoding
- Easy to extend to disjunctions and boolean combinations of DFAs

